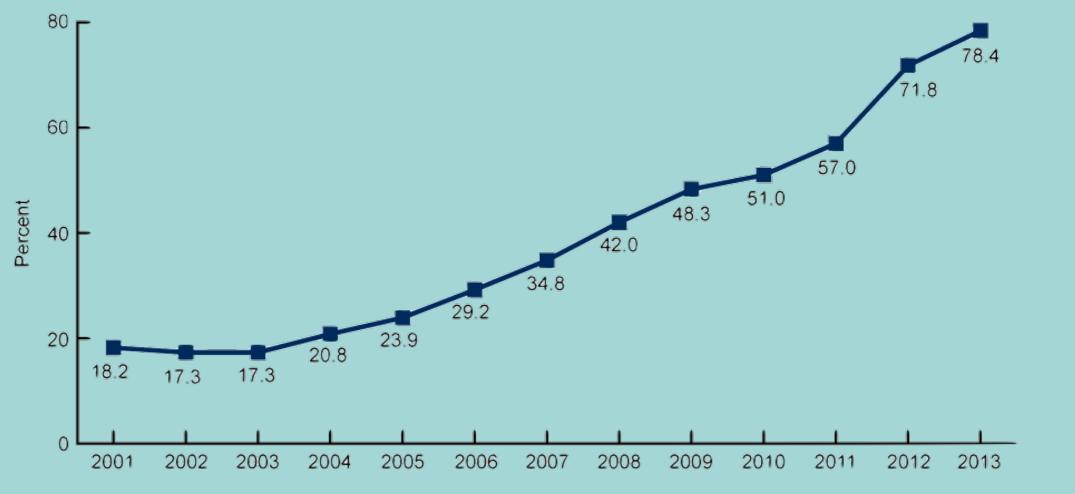
UNLYNX: A DECENTRALIZED SYSTEM FOR Privacy-Conscious Data Sharing PETS 2017

DAVID FROELICHER^{*#}, PATRICIA EGGER^{*#}, JOAO SA SOUSA^{*}, JEAN LOUIS RAISARO^{*}, Zhicong Huang^{*}, Christian Mouchet^{*}, Bryan Ford[#], Jean-Pierre Hubaux^{*}



MORE MEDICAL DATA ARE DIGITIZED

PERCENTAGE OF OFFICE-BASED PHYSICIANS WITH ELECTRONIC MEDICAL RECORDS IN U.S.A, 2001-2013



2

NATIONAL AMBULATORY MEDICAL CARE SURVEY (NAMCS)

MORE HEALTH DATA COLLECTED

3



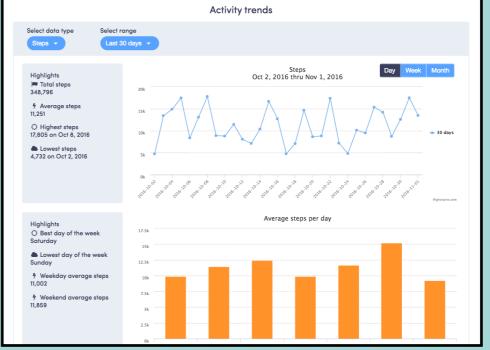
http://blog.stridekick.com/ultimate-guide-fitness-tracker-hacks-get-most-from-fitbit/



http://www.consumerreports.org/cro/news/2015/06/what-you-need-to-know-about-sharing-your-medical-data/index.htm



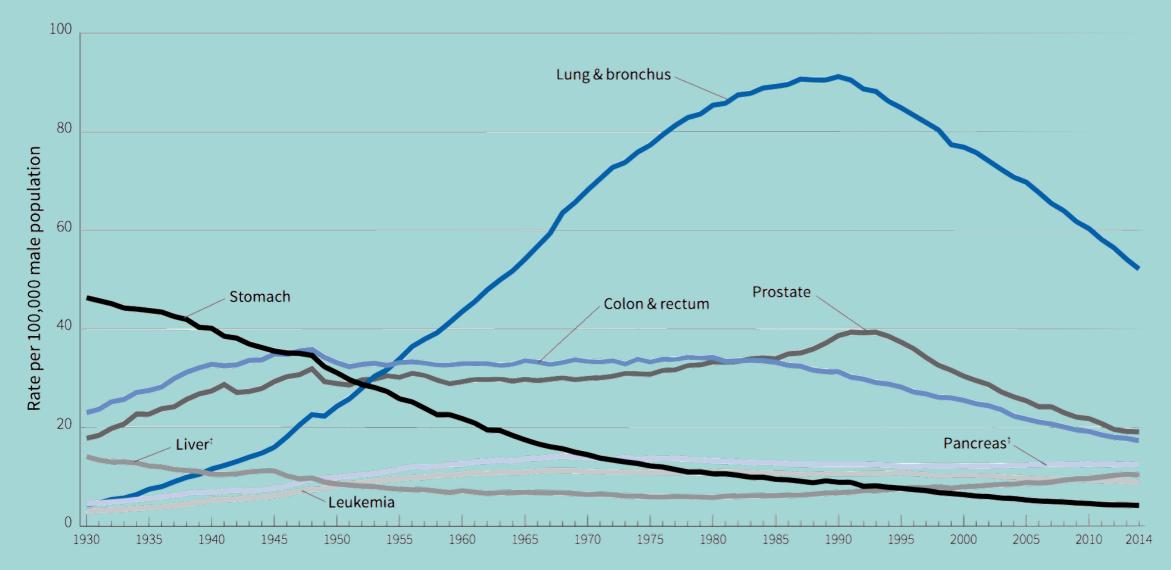
http://time.com/collection-post/3615161/sharing-health-data/



http://www.designindaba.com/articles/creative-work/smart-thermometer-crowdsources-info-real-time-health-tracking

MORE MEDICAL DATA = BETTER TREATMENTS ?

CANCER DEATH RATES* AMONG MEN, USA, 1930-2014



*Per 100,000, age adjusted to the 2000 US standard population. †Mortality rates for pancreatic and liver cancers are increasing.

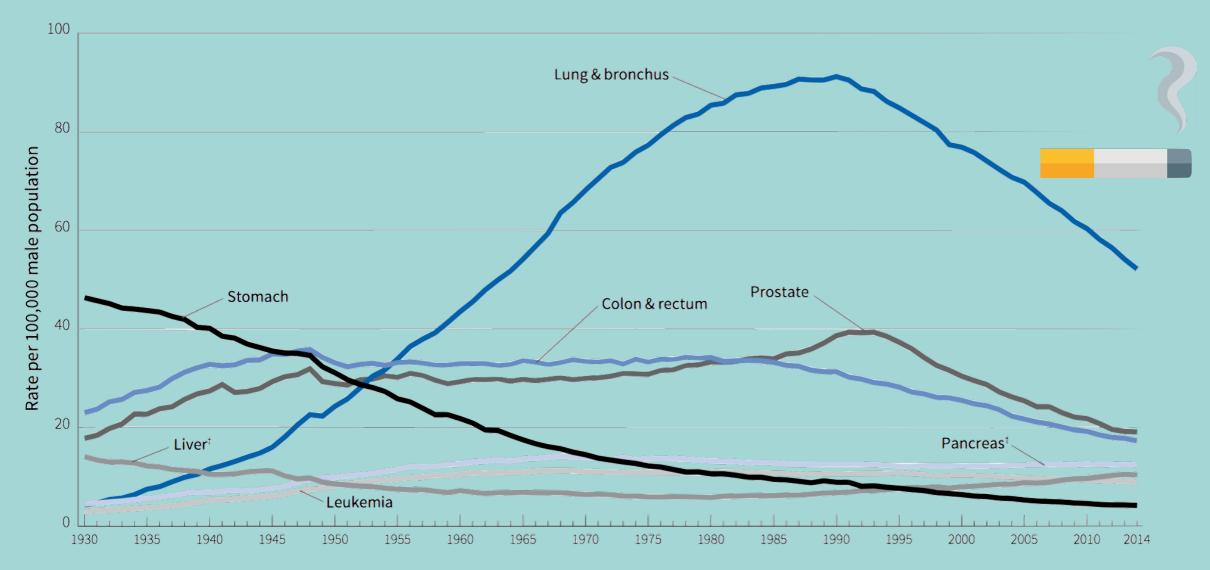
Note: Due to changes in ICD coding, numerator information has changed over time. Rates for cancers of the liver, lung and bronchus, uterus, and colon and rectum are affected by these coding changes.

Δ

Source: US Mortality Volumes 1930 to 1959 and US Mortality Data 1960 to 2014, National Center for Health Statistics, Centers for Disease Control and Prevention.

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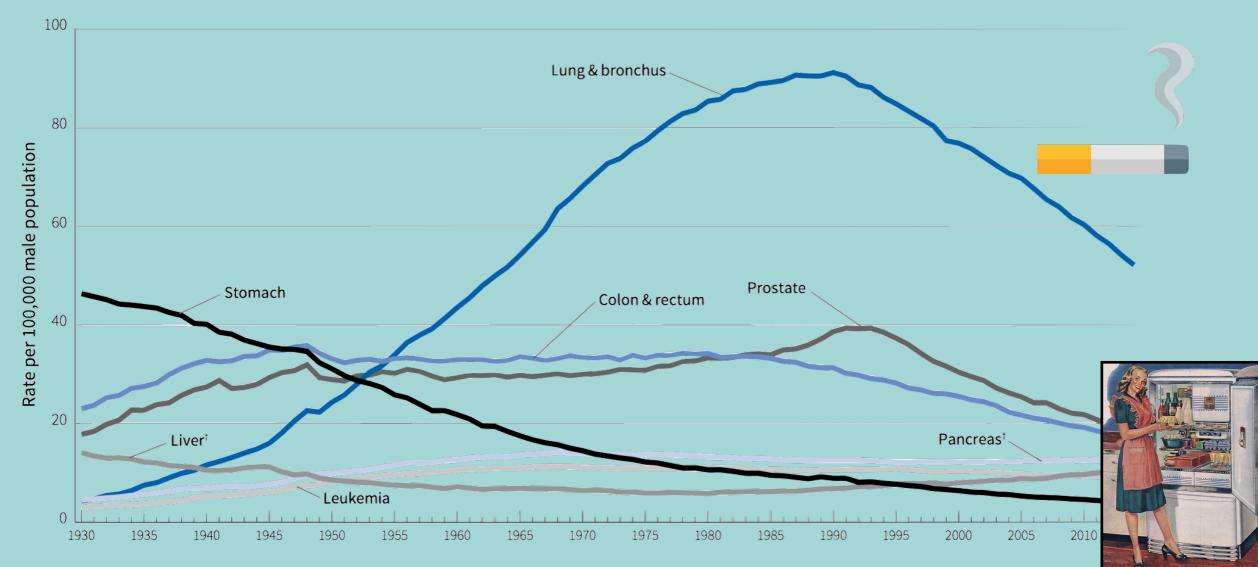
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5

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6

Source: US Mortality Volumes 1930 to 1959 and US Mortality Data 1960 to 2014, National Center for Health Statistics, Centers for Disease Control and Prevention.

SENSITIVE-DATA SHARING IS DIFFICULT



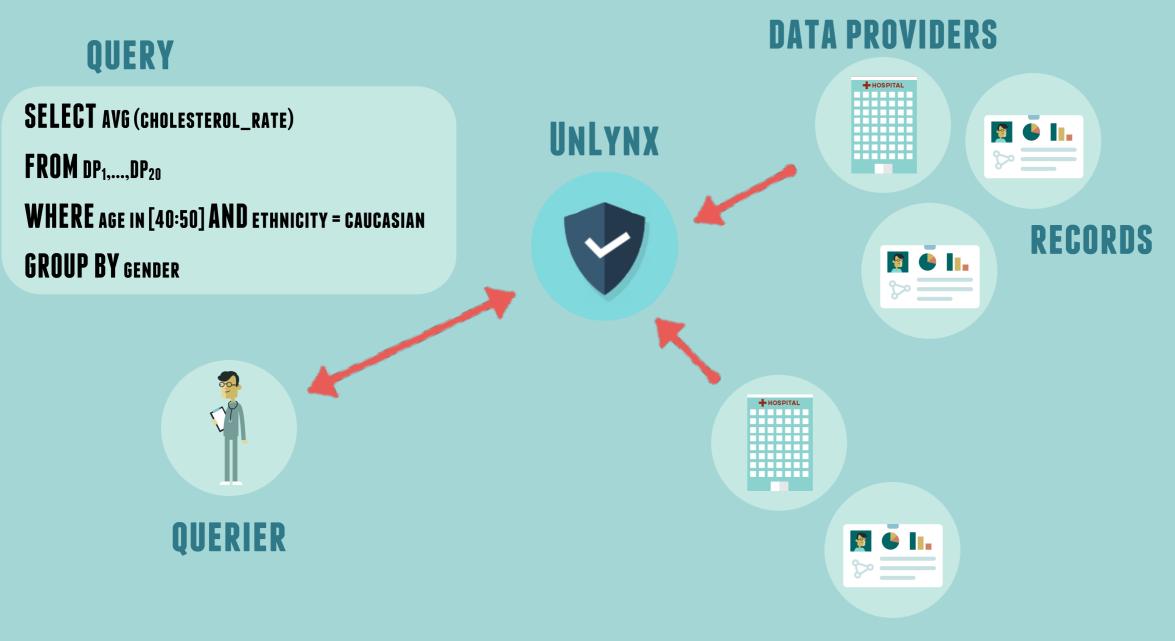
http://blog.lpinnovations.com

SENSITIVE-DATA SHARING IS DIFFICULT



http://www.gmmill.net/proje-Grain-Storage-Silos

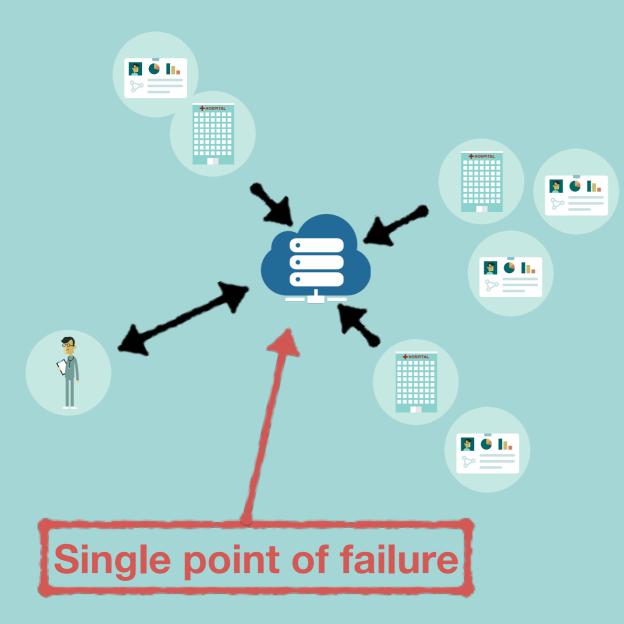
UNLYNX



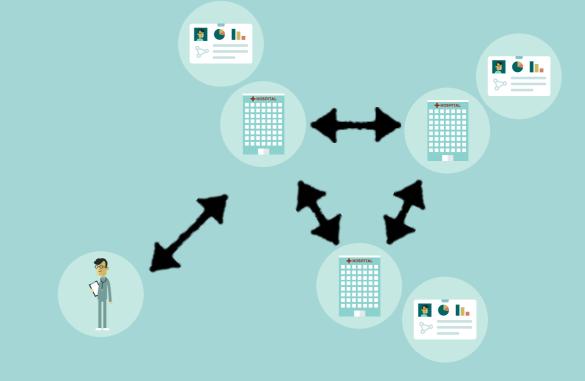
Allow statistical queries on multiple independent databases while ensuring privacy and confidentiality for data providers.

EXISTING DATA SHARING SOLUTIONS

CENTRALIZED SOLUTIONS



DECENTRALIZED SOLUTIONS



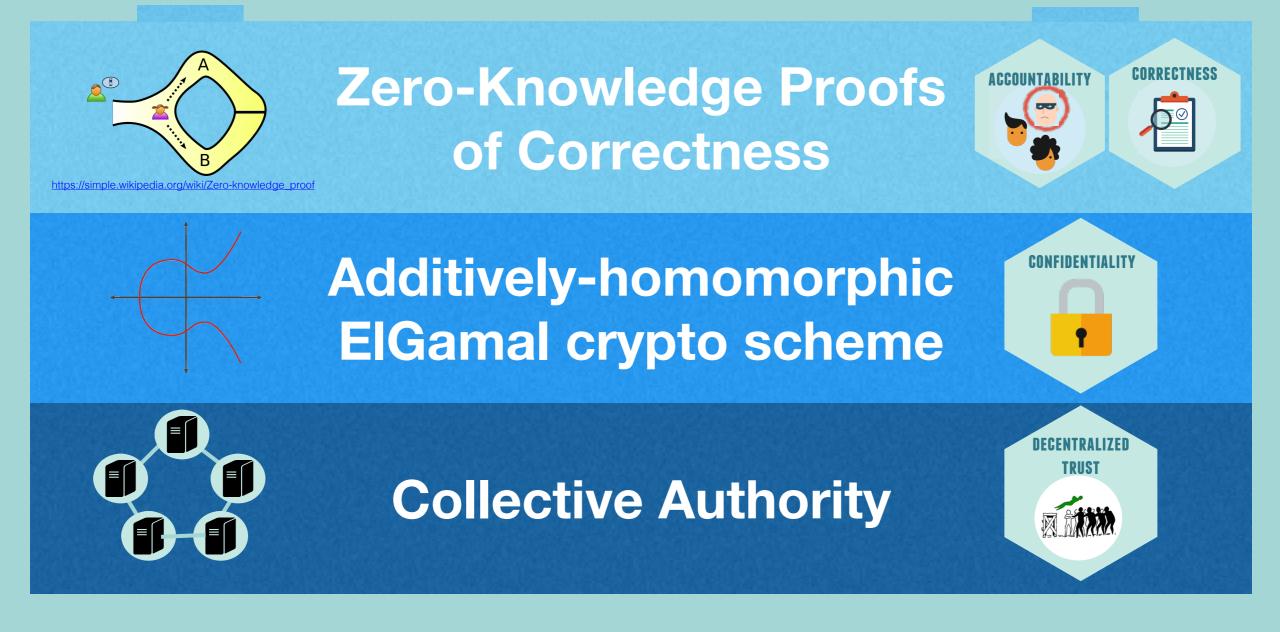
Limited number of data providers/computation entities in an adversarial model

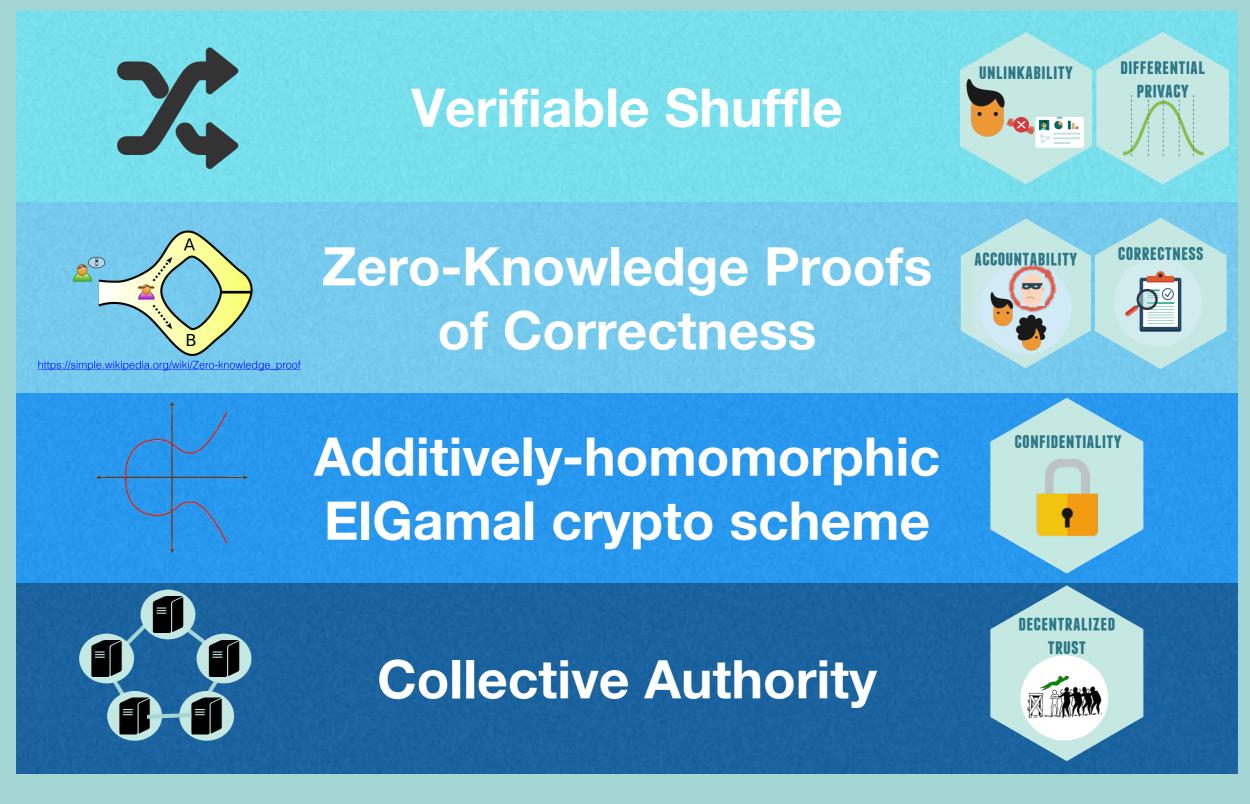
REQUIREMENTS



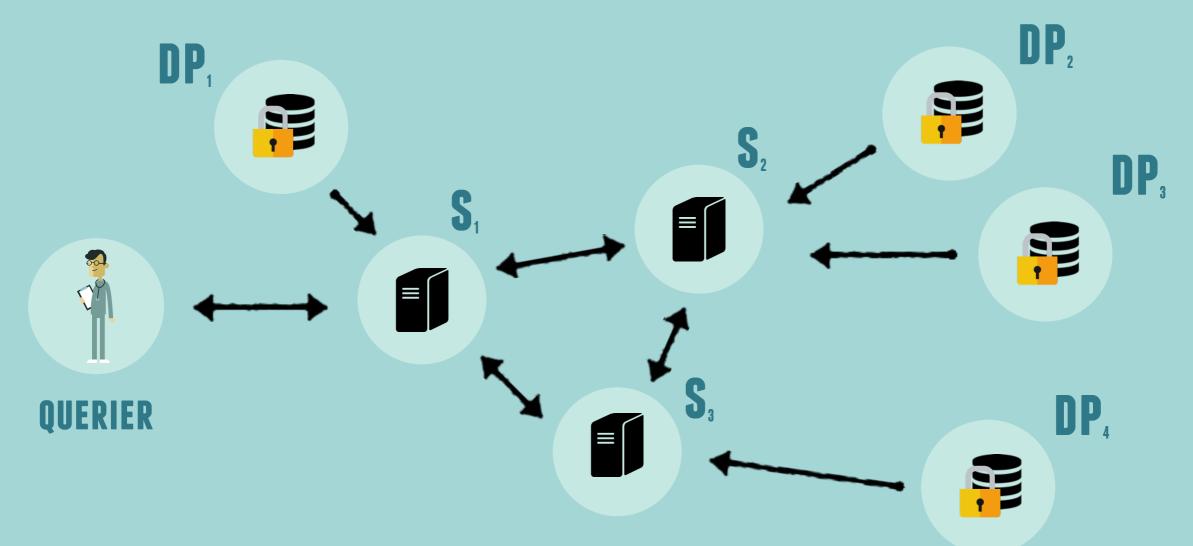








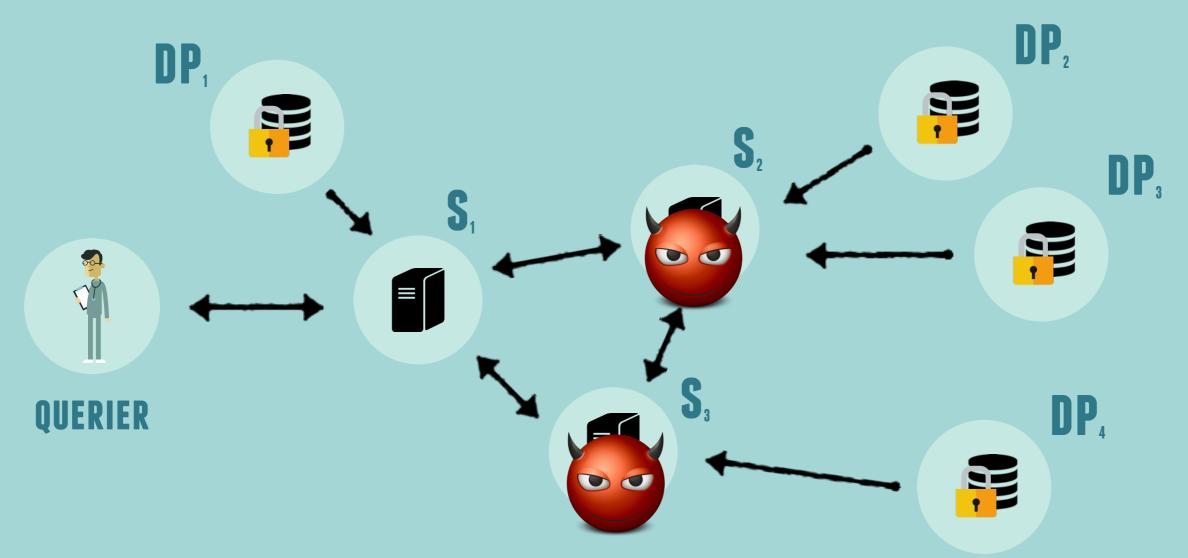
SYSTEM MODEL



- Collective authority of *m* servers S
- *n* Data Providers *DP*s
- Clients Q querying the system

DP = DATA PROVIDER S = Server

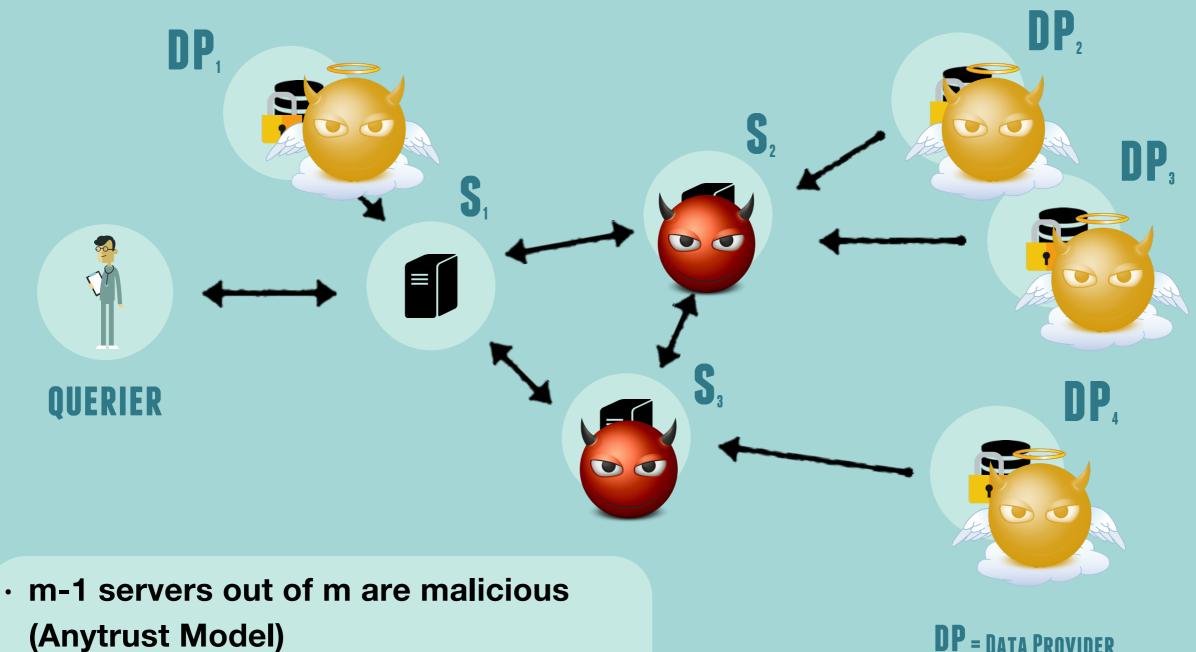
THREAT MODEL



 m-1 servers out of m are malicious (Anytrust Model)

DP = DATA PROVIDER S = Server

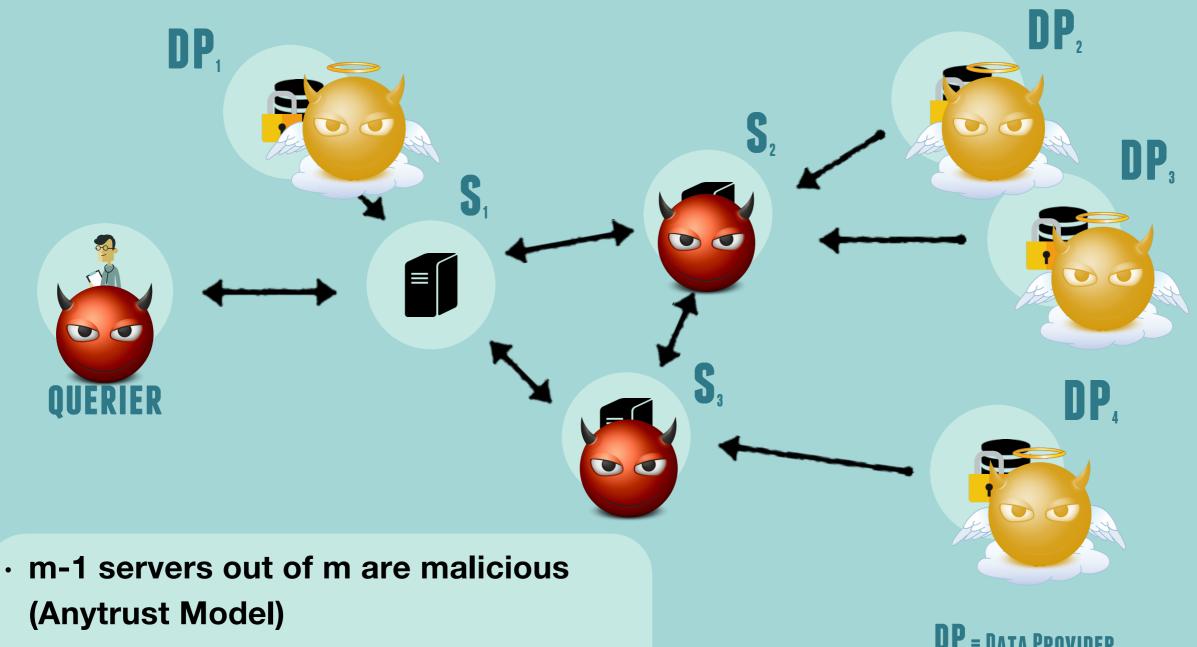
THREAT MODEL



Data Providers are honest-but-curious

DP = DATA PROVIDER S = SERVER

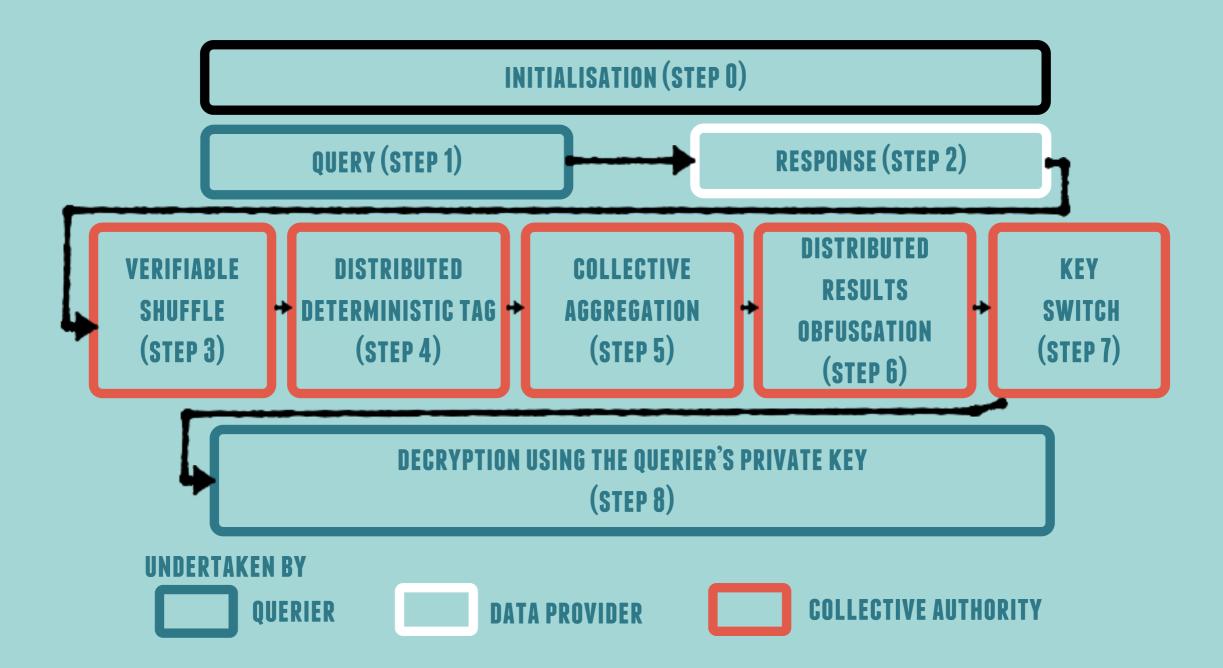
THREAT MODEL



- Data Providers are honest-but-curious
- Queriers are malicious

DP = DATA PROVIDER S = SERVER

QUERY PROCESSING WORKFLOW

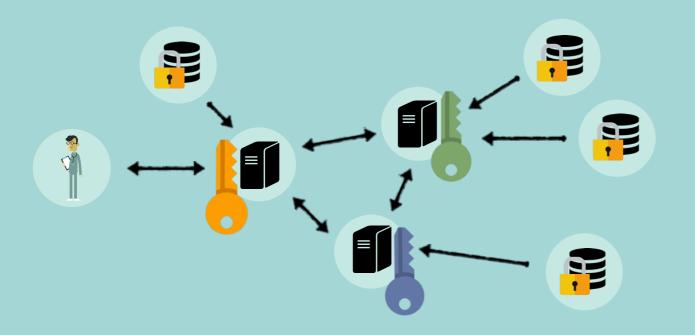


WORKFLOW - INITIALISATION (STEP 0)

INITIALISATION (STEP 0)



Each server constructs his publicprivate ElGamal Key pair.

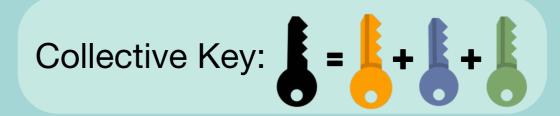


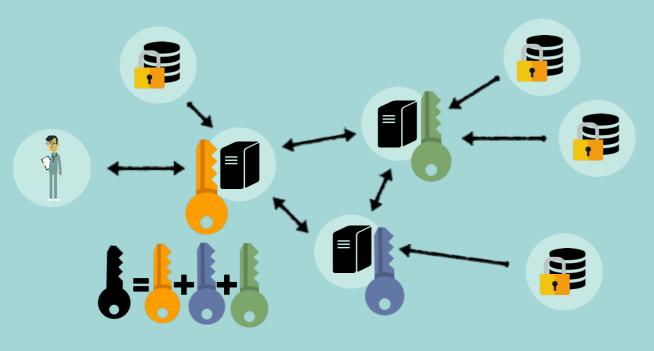
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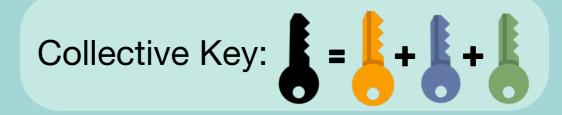


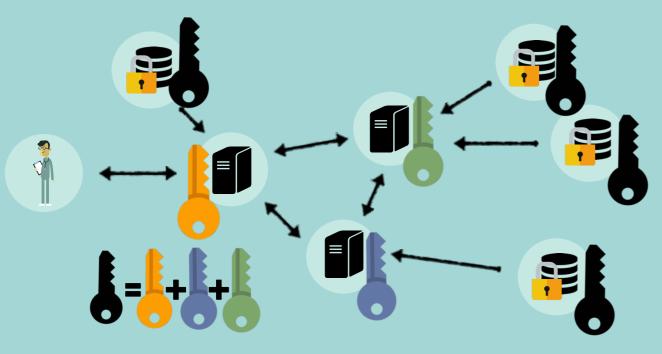
WORKFLOW - INITIALISATION (STEP 0)

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Each server constructs his publicprivate ElGamal Key pair.







Data Providers use the Collective Key to encrypt their data

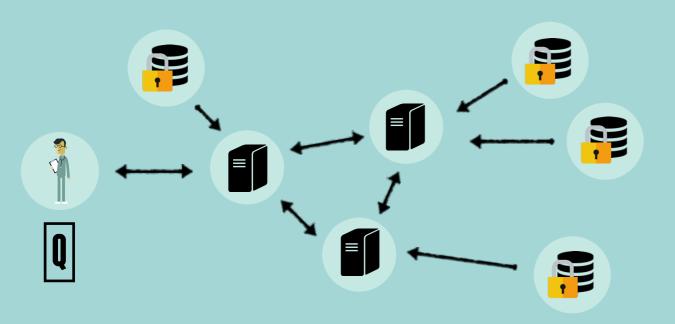
WORKFLOW - QUERY (STEP 1)

TN

INITIALISATION (STEP 0)

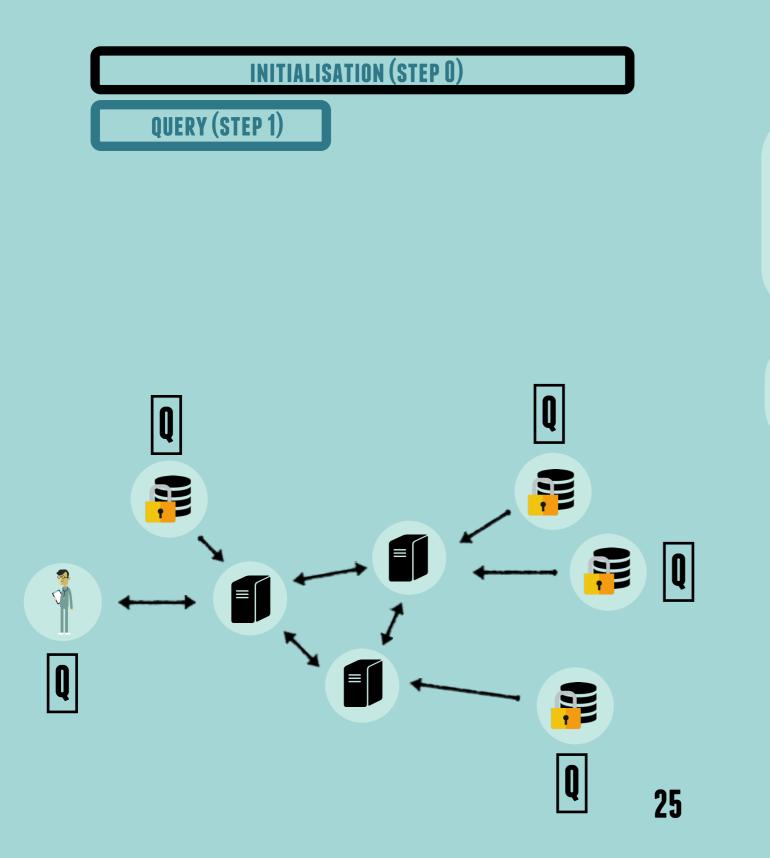


SELECT SUM (CHOLESTEROL_RATE), COUNT(*) FROM DP₁,...,DP₂₀ Where age in [40:50] and ethnicity = caucasian Group by gender



WORKFLOW - QUERY (STEP 1)

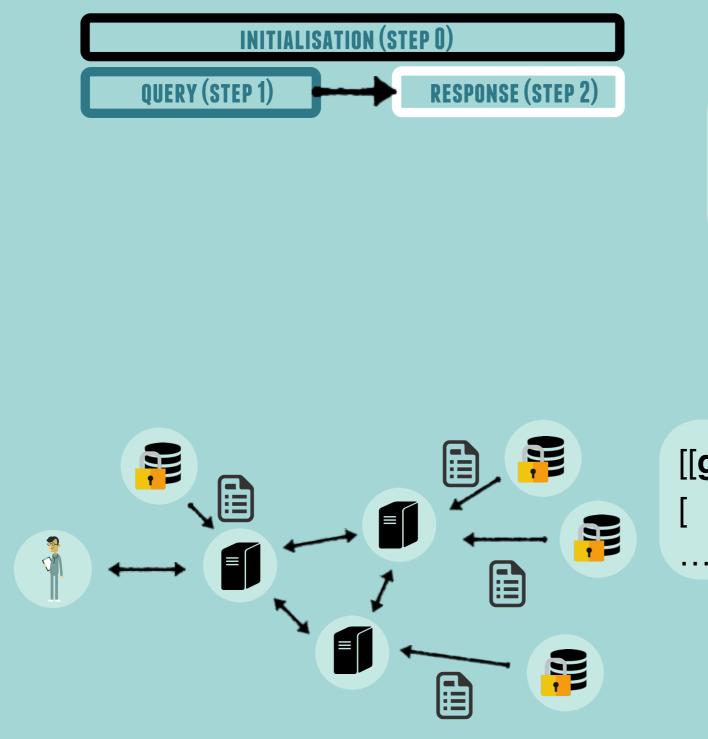
7



SELECT SUM (CHOLESTEROL_RATE), COUNT(*) FROM DP₁,...,DP₂₀ Where age in [40:50] AND ethnicity = caucasian GROUP BY gender

Query broadcasted to Data Providers

WORKFLOW - RESPONSE (STEP 2)

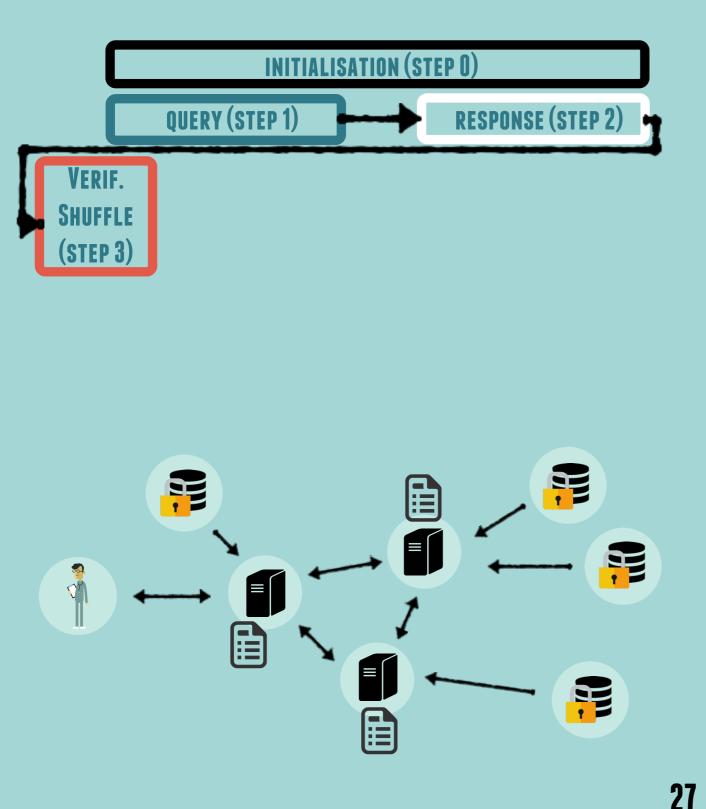


ID	Gender	Age	Ethnicity	flu	Cholesterol_rate	cancer
1	E ₁ (1)	E ₁ (40)	E _H (1)	E ₁ (1)	E ₁ (23)	E _E (0)
2	E (2)	E ₁ (40)	E ₁ (2)	E ₁ (0)	E.(34)	E.(0)



[[group. attr.], [where. attr.], [aggr. Attr.]] [$[E_{0}(1)]$, $[E_{0}(40), E_{0}(1)]$, $[E_{0}(23), E_{0}(0)]$

WORKFLOW - VERIF. SHUFFLE (STEP 3)



I

Each server starts a **verifiable shuffle protocol**:

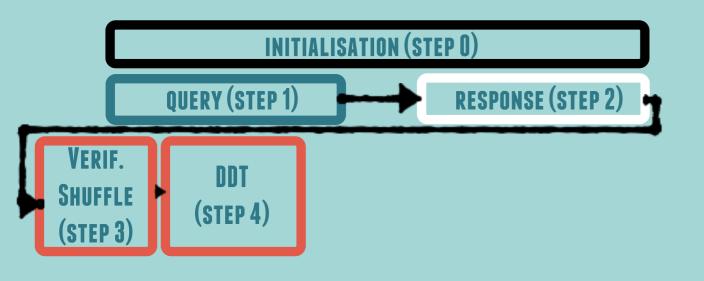
In this protocol each server sequentially:

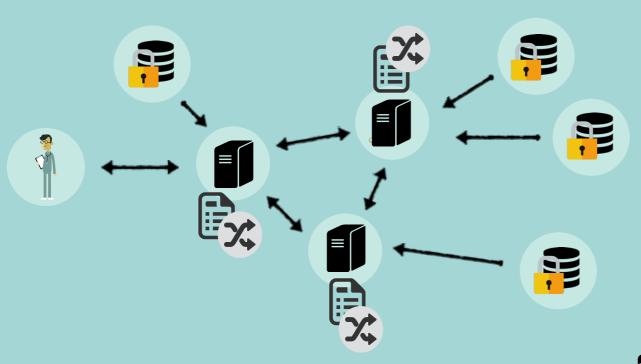
- Shuffle the list of responses
- Rerandomize (re-encryption) all the ciphertexts

Using Neff Shuffle and the corresponding zero-knowledge proof [1]

[1] Andrew Neff. Verifiable mixing (shuffling) of ElGamal pairs (2004)

WORKFLOW - DDT (STEP 4)



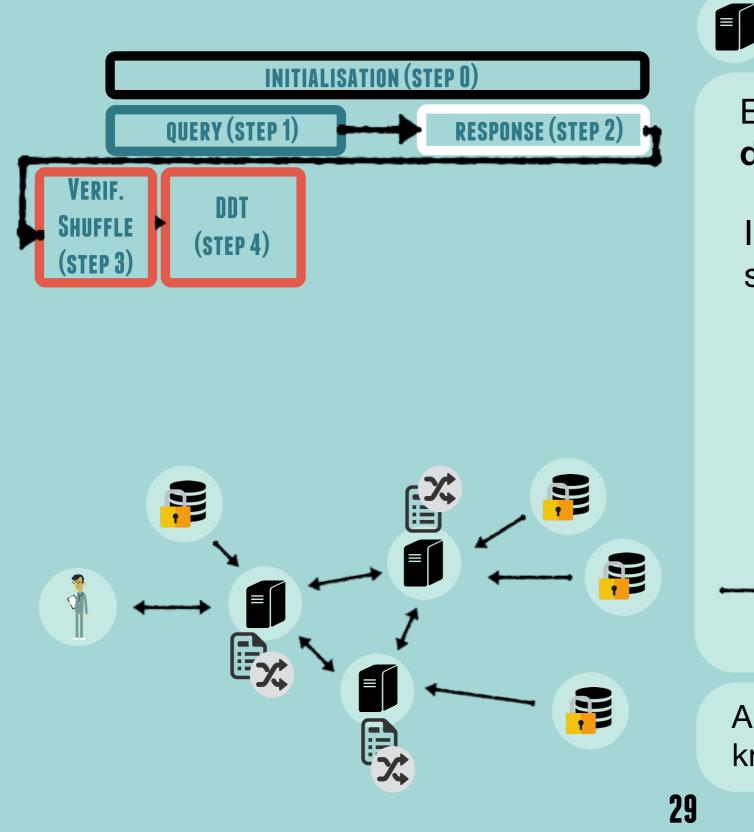


Each server starts a **distributed deterministic tagging protocol**:

Query: WHERE $age = E\kappa(40)$ AND $ethnicity = E\kappa(2)$ WHERE age = DT(40) AND ethnicity = DT(2)

Data: [[Ек(1)], [Ек(40),Ек(2)], [Ек(23),Ек(1)] ↓ [[DT(1)], [DT(40),DT(2)], [Ек(23),Ек(1)]

WORKFLOW - DDT (STEP 4)



Each server starts a **distributed deterministic tagging protocol**:

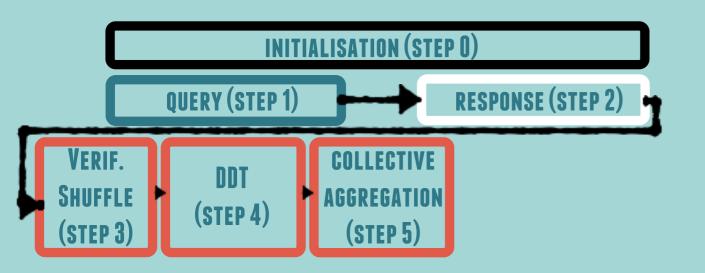
In this protocol each server sequentially:

- partially decrypt the ciphertexts
- Blinds the message by multiplying the ciphertexts with a random ephemeral secret key
- deterministic tag depending on the value of the encrypted message

All operations are done with zeroknowledge proofs from Camenisch et al.

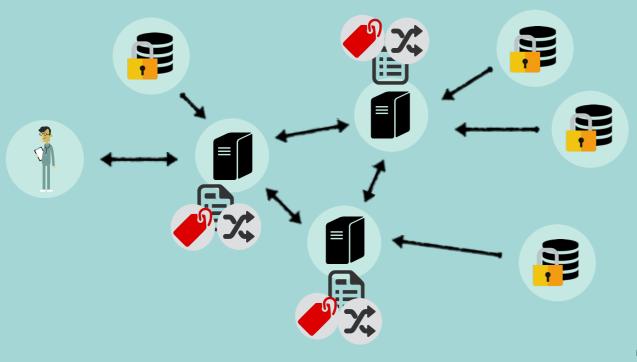
[1] Jan Camenish and Markus Stadler. Proof systems for general statements about discrete logarithms. (1997)

WORKFLOW - COLLECTIVE AGGR. (STEP 5)



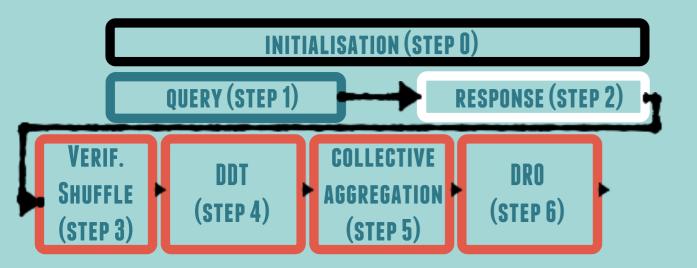


Servers **collectively aggregate** the responses by group.



Proofs consist in publishing the ciphertexts and the result

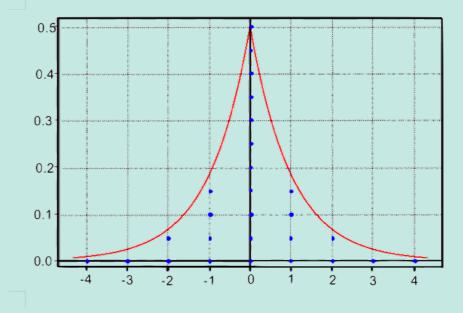
WORKFLOW - DRO (STEP 6)



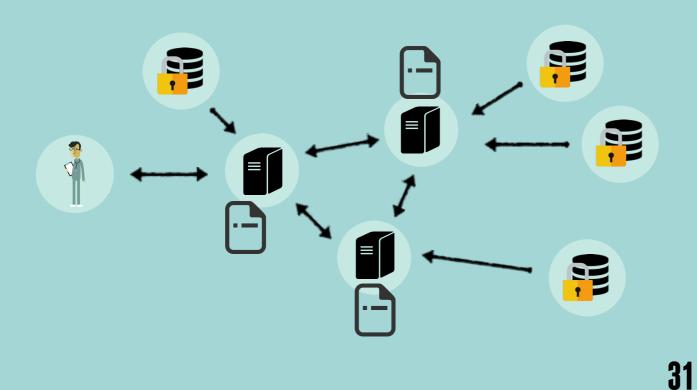


Distributed Results Obfuscation: <u>Setup:</u>

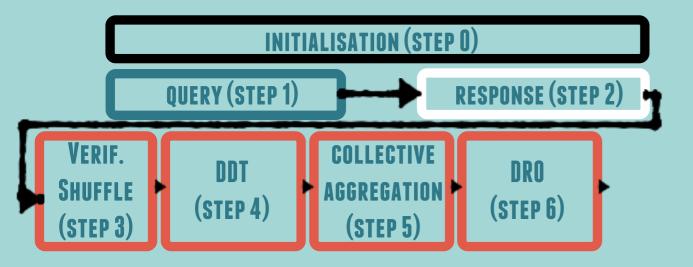
Servers agree on (ϵ, δ) -differential privacy parameters and produce:

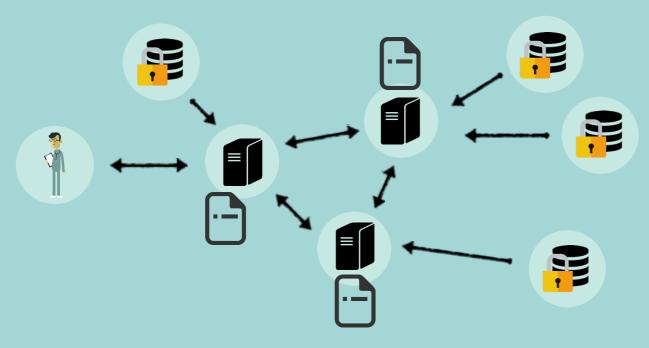


 $\longrightarrow [0,0,0,0,0,0,0,0,0,0,0,1,1,1,1,1,1,...] =$ list of noise values satisfying (ε,δ)- differential privacy.



WORKFLOW - DRO (STEP 6)



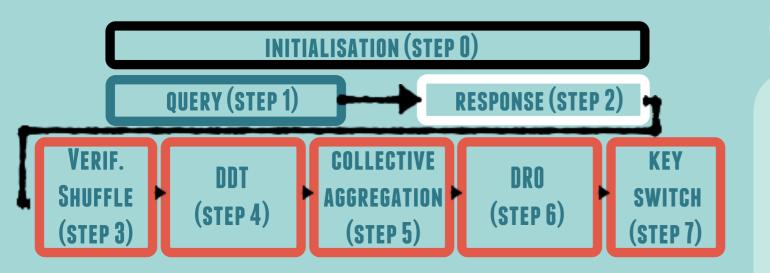




Distributed Results Obfuscation: <u>Runtime:</u>

- A server starts a collective shuffling of the list of noise values
- adds the first noise value in the list to the query result.
- → Oblivious noise addition (shuffling encrypts and shuffles the list of noise values).

WORKFLOW - KEY SWITCH (STEP 7)



I

In the **key switch protocol** each server:

- partially decrypt
- **encrypt** with a new key all the ciphertexts.

→ Encryption is switched from the Collective Key to the querier's public key.

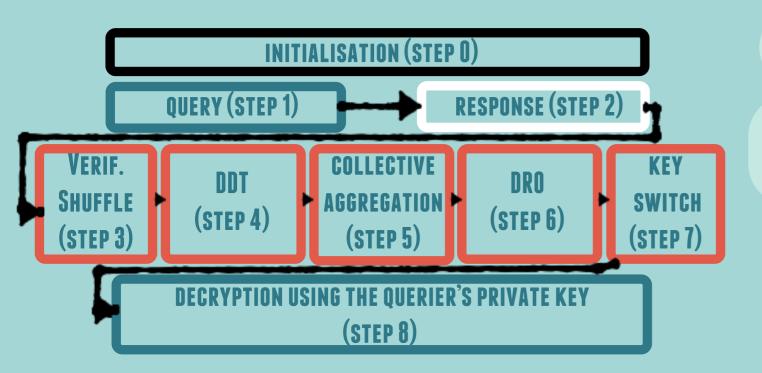


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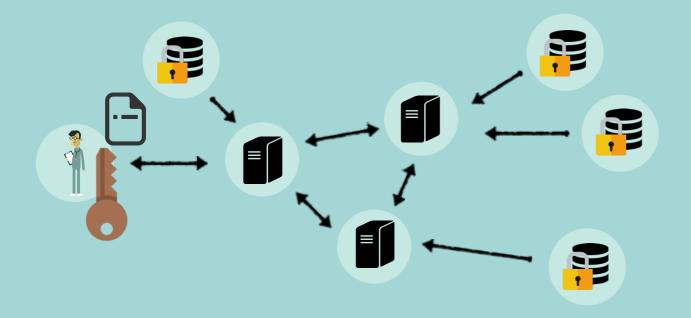
[1] Jan Camenish and Markus Stadler. Proof systems for general statements about discrete logarithms. (1997)

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WORKFLOW - DECRYPTION (STEP 8)



Querier **decrypts** the result with his secret key



PERFORMANCE EVALUATION

Servers configuration

- Memory: 256GB RAM
- Processor: Intel Xeon E5-2680 v3 (Haswell)
- Cores: 24 (with 48 threads)
- Frequency: 2.5GHz
- Bandwidth capacity: 1Gbps

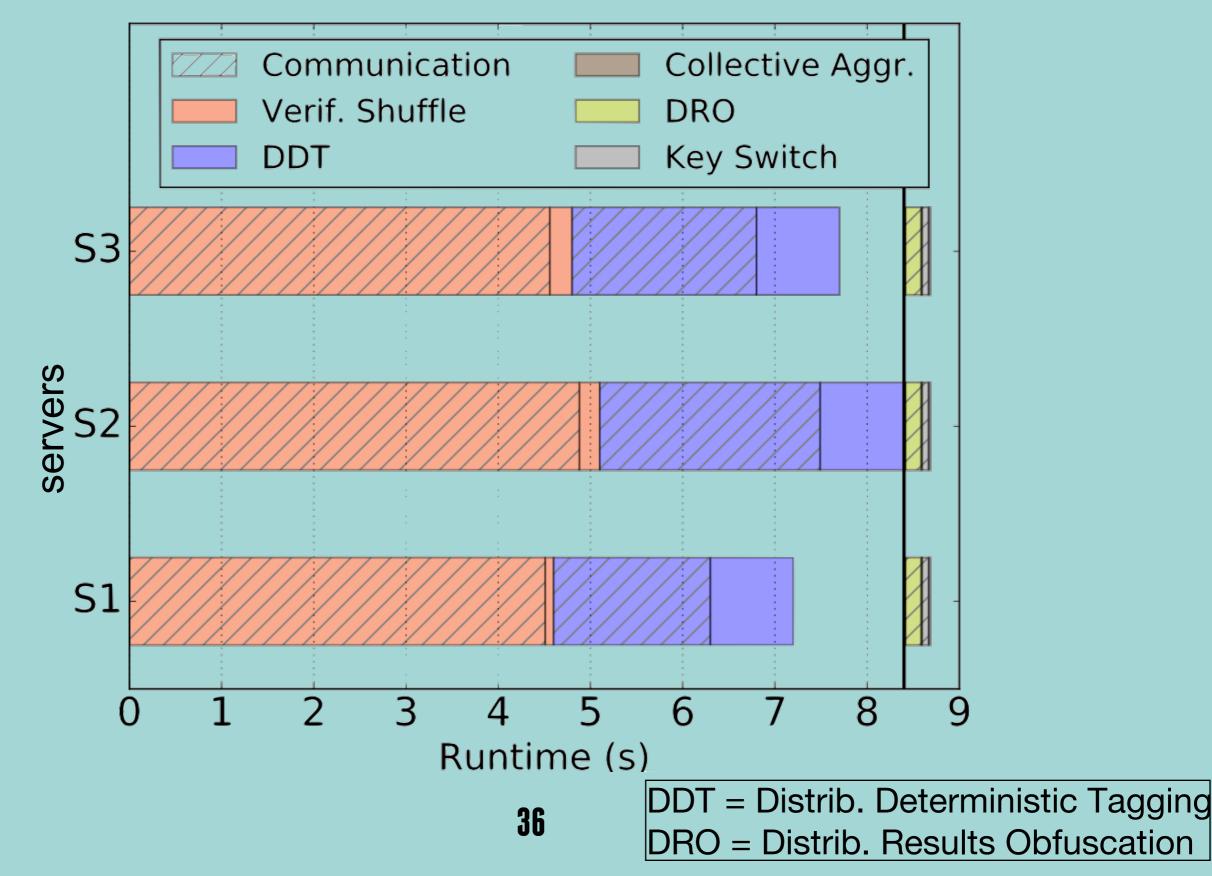
Network and Crypto

- Realistic virtual network emulation tool with 10ms delays btw. servers
- DeDiS' Onet library
- DeDiS' implementation of Ed25519 Elliptic Curve (128-bit security)

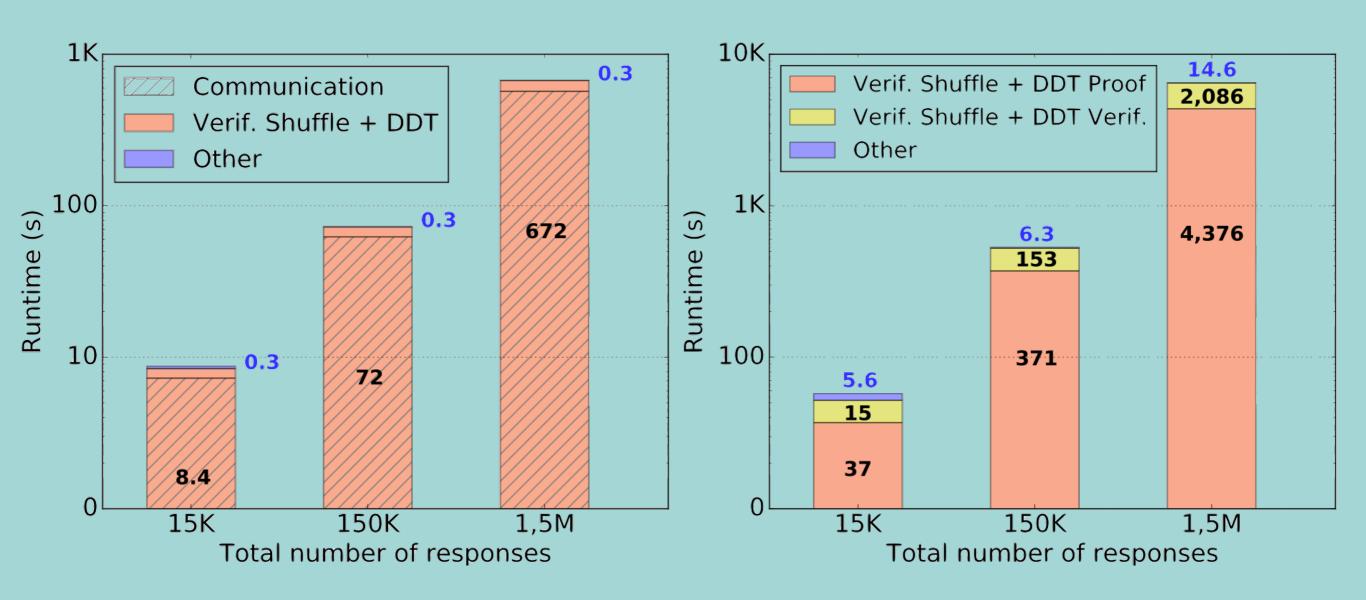
Default parameters

- 3 servers
- 15,000 responses in total (equally distributed in servers)
- 1 GROUP BY attribute with 10 possible values , 1 WHERE and 10 aggregating attributes
- 1000 noise values

SERVERS COLLABORATION



RUNTIME VS. NBR. OF RESPONSES



PERFORMANCE/SECURITY TRADEOFFS



SELECT SUM (CHOLESTEROL_RATE), COUNT(*) FROM DP1,...,DP20 WHERE AGE IN [40:50] AND ETHNICITY = CAUCASIAN GROUP BY GENDER

3 servers 400K responses with 1 GROUP BY attribute 2 WHERE attributes 2 aggregating attributes

CONCLUSION

A Decentralized System for Privacy-Conscious Data Sharing

- SQL statistical queries based on Boolean conditions
- Strongest-link security
- Data confidentiality
- Distributed differential privacy
- Distributed deterministic tagging of probabilistic ciphertexts
- Collective encryption key switching
- Runtime linear with the amount of data to process

github.com/lca1/unlynx















david.froelicher@epfl.ch